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APPLICATION NO.	FILING DATE	FIRST NAMED INVENTOR	ATTORNEY DOCKET NO.	CONFIRMATION NO.
10/648,311	08/27/2003	Sathya P. Kaginele	M4065.0931/P931	4875
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DICKSTEIN SHAPIRO LLP 1825 EYE STREET, NW WASHINGTON, DC 20006			EXAMINER TU, CHRISTINE TRINH LE	
			ART UNIT 2117	PAPER NUMBER
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Please find below and/or attached an Office communication concerning this application or proceeding.

The time period for reply, if any, is set in the attached communication.

Office Action Summary	Application No. 10/648,311	Applicant(s) KAGINELE, SATHYA P.	
	Examiner Christine T. Tu	Art Unit 2117	

-- The MAILING DATE of this communication appears on the cover sheet with the correspondence address --

Period for Reply

A SHORTENED STATUTORY PERIOD FOR REPLY IS SET TO EXPIRE 3 MONTH(S) OR THIRTY (30) DAYS, WHICHEVER IS LONGER, FROM THE MAILING DATE OF THIS COMMUNICATION.

- Extensions of time may be available under the provisions of 37 CFR 1.136(a). In no event, however, may a reply be timely filed after SIX (6) MONTHS from the mailing date of this communication.
- If NO period for reply is specified above, the maximum statutory period will apply and will expire SIX (6) MONTHS from the mailing date of this communication.
- Failure to reply within the set or extended period for reply will, by statute, cause the application to become ABANDONED (35 U.S.C. § 133). Any reply received by the Office later than three months after the mailing date of this communication, even if timely filed, may reduce any earned patent term adjustment. See 37 CFR 1.704(b).

Status

- 1) ☒ Responsive to communication(s) filed on 4/19/2007.
- 2a) ☒ This action is **FINAL**. 2b) ☐ This action is non-final.
- 3) ☐ Since this application is in condition for allowance except for formal matters, prosecution as to the merits is closed in accordance with the practice under *Ex parte Quayle*, 1935 C.D. 11, 453 O.G. 213.

Disposition of Claims

- 4) ☒ Claim(s) 1-10,12-14,17,19-23,30,32-36 and 38-42 is/are pending in the application.
- 4a) Of the above claim(s) _____ is/are withdrawn from consideration.
- 5) ☐ Claim(s) _____ is/are allowed.
- 6) ☒ Claim(s) 1-10,12-14,17,19-23,30,32-36 and 38-42 is/are rejected.
- 7) ☐ Claim(s) _____ is/are objected to.
- 8) ☐ Claim(s) _____ are subject to restriction and/or election requirement.

Application Papers

- 9) ☐ The specification is objected to by the Examiner.
- 10) ☐ The drawing(s) filed on _____ is/are: a) ☐ accepted or b) ☐ objected to by the Examiner.
Applicant may not request that any objection to the drawing(s) be held in abeyance. See 37 CFR 1.85(a).
Replacement drawing sheet(s) including the correction is required if the drawing(s) is objected to. See 37 CFR 1.121(d).
- 11) ☐ The oath or declaration is objected to by the Examiner. Note the attached Office Action or form PTO-152.

Priority under 35 U.S.C. § 119

- 12) ☐ Acknowledgment is made of a claim for foreign priority under 35 U.S.C. § 119(a)-(d) or (f).
- a) ☐ All b) ☐ Some * c) ☐ None of:
1. ☐ Certified copies of the priority documents have been received.
 2. ☐ Certified copies of the priority documents have been received in Application No. _____.
 3. ☐ Copies of the certified copies of the priority documents have been received in this National Stage application from the International Bureau (PCT Rule 17.2(a)).

* See the attached detailed Office action for a list of the certified copies not received.

Attachment(s)

- | | |
|--|---|
| 1) <input type="checkbox"/> Notice of References Cited (PTO-892) | 4) <input type="checkbox"/> Interview Summary (PTO-413)
Paper No(s)/Mail Date. _____ |
| 2) <input type="checkbox"/> Notice of Draftsperson's Patent Drawing Review (PTO-948) | 5) <input type="checkbox"/> Notice of Informal Patent Application |
| 3) <input type="checkbox"/> Information Disclosure Statement(s) (PTO/SB/08)
Paper No(s)/Mail Date _____ | 6) <input type="checkbox"/> Other: _____ |

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1. The text of those sections of Title 35, U.S. Code not included in this action can be found in a prior Office action.

Double Patenting

2. The nonstatutory double patenting rejection is based on a judicially created doctrine grounded in public policy (a policy reflected in the statute) so as to prevent the unjustified or improper timewise extension of the "right to exclude" granted by a patent and to prevent possible harassment by multiple assignees. A nonstatutory obviousness-type double patenting rejection is appropriate where the conflicting claims are not identical, but at least one examined application claim is not patentably distinct from the reference claim(s) because the examined application claim is either anticipated by, or would have been obvious over, the reference claim(s). See, e.g., *In re Berg*, 140 F.3d 1428, 46 USPQ2d 1226 (Fed. Cir. 1998); *In re Goodman*, 11 F.3d 1046, 29 USPQ2d 2010 (Fed. Cir. 1993); *In re Longi*, 759 F.2d 887, 225 USPQ 645 (Fed. Cir. 1985); *In re Van Ornum*, 686 F.2d 937, 214 USPQ 761 (CCPA 1982); *In re Vogel*, 422 F.2d 438, 164 USPQ 619 (CCPA 1970); and *In re Thorington*, 418 F.2d 528, 163 USPQ 644 (CCPA 1969).

A timely filed terminal disclaimer in compliance with 37 CFR 1.321(c) or 1.321(d) may be used to overcome an actual or provisional rejection based on a nonstatutory double patenting ground provided the conflicting application or patent either is shown to be commonly owned with this application, or claims an invention made as a result of activities undertaken within the scope of a joint research agreement.

Effective January 1, 1994, a registered attorney or agent of record may sign a terminal disclaimer. A terminal disclaimer signed by the assignee must fully comply with 37 CFR 3.73(b).

3. Claims (1 & 2), 3, 6, 7 are provisionally rejected on the ground of nonstatutory obviousness-type double patenting as being unpatentable over claims (47 & 48), 49, 50, 51 of copending Application No. 11/474,495, respectively. Although the conflicting claims are not identical, they are not patentably distinct from each other because the copending application substantially recites the claimed invention. The copending application '495 does not explicitly disclose the feature of confirming proper operation of a control line used to enable output from a match line under test. It would have been

obvious to one skilled in the art at the time the invention was made to realize a feature of confirming proper operation of a control line would have been necessary. One having ordinary skill in the art would be motivated to realize so because the copending application '495 discloses the feature of enabling output from a match line (in claim 47).

This is a provisional obviousness-type double patenting rejection because the conflicting claims have not in fact been patented.

4. Claims (1 & 42, 2, 4, & 5), 6 and 7 are provisionally rejected on the ground of nonstatutory obviousness-type double patenting as being unpatentable over claims 52, 53 and 54 of copending Application No. 11/474,495, respectively. Although the conflicting claims are not identical, they are not patentably distinct from each other because the copending application substantially recites the claimed invention. The copending application '495 does not explicitly disclose the feature of confirming proper operation of a control line used to enable output from a match line under test. It would have been obvious to one skilled in the art at the time the invention was made to realize a feature of confirming proper operation of a control line would have been necessary. One having ordinary skill in the art would be motivated to realize so because the copending application '495 discloses the feature of enabling match line (in claim 52.

This is a provisional obviousness-type double patenting rejection because the conflicting claims have not in fact been patented.

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5. Claims 8, 10, 13 and 19 are provisionally rejected on the ground of nonstatutory obviousness-type double patenting as being unpatentable over claims 43, 44, 45 and 46 of copending Application No. 11/474,495, respectively. Although the conflicting claims are not identical, they are not patentably distinct from each other because the copending application substantially recites the claimed invention. The copending application '495 does not explicitly disclose the feature of confirming proper operation of a control line used to enable output from a match line under test. It would have been obvious to one skilled in the art at the time the invention was made to realize a feature of confirming proper operation of a control line would have been necessary. One having ordinary skill in the art would be motivated to realize so because the copending application '495 discloses the feature of enabling output from a match line (in claim 43).

This is a provisional obviousness-type double patenting rejection because the conflicting claims have not in fact been patented.

Claim Rejections - 35 USC § 103

6. Claims 1-10, 12-14, 17, 19-23, 30, 32-36, and 38-42 are rejected under 35 U.S.C. 103(a) as being unpatentable over Ichiriu (7,002,823) in view of Nataraj et al. (6,944,039 and Nataraj hereinafter).

Claims 1-2:

Ichiriu discloses the invention substantially as claimed. Ichiriu shows (figures 1 & 5) that, in a CAM device (100), a CAM array (101) includes a plurality of CAM cells arranged in rows for storing CAM words. Each CAM cell (201) includes a compare circuit to compare the content of the memory cell with a comparand signal. A comparand register (115) is used to store a comparand value received via the comparand bus (143) and outputs the comparand value to the CAM array (101). During a compare operation, the comparand may be masked by a global mask value, and then compared simultaneously with all the CAM words stored in the CAM array. Each of the rows of CAM cell is coupled to a corresponding match line (182) and any match between the comparand and a valid CAM word results in a match signal being output to the flag circuit (112). The flag circuit (112) then outputs a match flag signal to indicate that a match has occurred (figures 1 and 5, column 3 line 29 - column 4 line 34, column 6 lines 42-47).

Ichiriu states that a decoder (103) inside the CAM device (100) decodes a selected address to activate one of a plurality of word lines (181) (column 5 lines 17-23, column 3 lines 54-56).

Ichiriu also states at the start of a comparison operation, a feature of precharging each of the match lines (182) to a high logical level (column 6 line 61-column 7 line 4).

Ichiriu further states that, during a compare operation, a respective portion of the comparand is applied to each column of CAM cells (201) via lines CL/CLB such that the complete comparand is applied to each row of the CAM cells (201) simultaneously. In one embodiment, each of the match lines (182) is precharged to a high logical level at the start of a comparison operation, but pulled down to a low logical level by the compare circuit with any attached CAM cell (201) that receives comparand signals which do not match the stored data value. In this configuration, any match line (182) NOT pull low constitutes a match signal (figure 5, column 6 line 61-column 7 line 4).

Ichiriu does not explicitly teach the feature of confirming proper operation of a control line used to enable output from a match line under test and the feature of enabling output from the match line under test.

Nataraj, however, teaches (figure 4) that a match latch circuit (203) comprises latching circuits (221), each latching circuit comprises an AND gate (223) and a latch element (225). Each AND gate (223) has a first input coupled to receive a detect signal (224) and a second input coupled to a corresponding one of the match lines (241). The AND logic gate (223) outputs a logic-level match signal in high or low state according to whether the signal level of the corresponding match line is above or below a threshold. The latch signal (226) is supplied to a latch enable input of the latch element (225) enables the logic-level match signal to pass through the latch element (225) (figure 4, column 6 lines 48-62).

It would have been obvious to one skilled in the art at the time the invention was made to use Nataraj's match latch circuit (203) having the plurality of AND gates (223) and plurality of latch elements (225) such that each latch element (225) is enabled to pass/output a match signal [generated by a AND gate (223)] to Ichiriu's priority encoder (114) and Ichiriu's FLAG circuit (112). One having ordinary skill in the art would be motivated to combine the teachings of Ichiriu and Nataraj because the combination of Ichiriu's priority encoder [114] & FLAG circuit [112] and Nataraj's priority encoder/flag logic circuit (215) comprises inputs of match signals and outputs of match addresses and match flag (Ichiriu: figure 1, column 3 line 64-column 4 line 13; Nataraj: figure 4, column 7 line 1-15).

Claim 3:

Ichiriu teaches that the CAM array (101) includes circuitry to force validity value with each validity storage cell (202) to a reset state to prevent assertion of a match signal by pulling the match line low for the corresponding row of CAM cells (column 7 lines 49-65, column 8 lines 20-24).

Claims 4-5:

Ichiriu's decoder (105) is a state machine that transitions from state to state in response to transitions of a clock signal (CLK) (104) (column 4 lines 57-67).

Claims 6-7:

Ichiriu teaches that each row of CAM cells (201) is coupled to a respective match line (182) (figure 5, column 6 lines 33-37).

Claims 8-9:

Ichiriu discloses the invention substantially as claimed. Ichiriu shows (figure 1) that a CAM array (101) includes a plurality of CAM cell arranged in rows for storing CAM words. Each CAM cell (201) includes a compare circuit to compare the content of the memory cell with a comparand signal. A comparand register (115) is used to store a comparand value received via the comparand bus (143) and outputs the comparand value to the CAM array (101). During a compare operation, the comparand may be masked by a global mask value, and then compared simultaneously with all the CAM words stored in the CAM array. Each of the rows of CAM cell is coupled to a corresponding match line (182) and any match between the comparand and a valid CAM word results in a match signal being output to the flag circuit (112). The flag circuit (112) then outputs a match flag signal to indicate that a match has occurred (figure 1, column 3 line 29 - column 4 line 34, column 6 lines 42-47).

Ichiriu does not explicitly teach the feature of confirming proper operation of a control line used to enable output from a match line under test and the feature of enabling output from the match line under test.

Nataraj, however, teaches (figure 4) that a match latch circuit (203) comprises latching circuits (221), each latching circuit comprises an AND gate (223) and a latch element (225). Each AND gate (223) has a first input coupled to receive a detect signal

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(224) and a second input coupled to a corresponding one of the match lines (241). The AND logic gate (223) outputs a logic-level match signal in high or low state according to whether the signal level of the corresponding match line is above or below a threshold. The latch signal (226) is supplied to a latch enable input of the latch element (225) enables the logic-level match signal to pass through the latch element (225) (figure 4, column 48-62).

It would have been obvious to one skilled in the art at the time the invention was made to use Nataraj's match latch circuit (203) having the plurality of AND gates (223) and plurality of latch elements (225) such that each latch element (225) is enabled to pass/output a match signal [generated by a AND gate (223)] to Ichiriu's priority encoder (114) and Ichiriu's FLAG circuit (112). One having ordinary skill in the art would be motivated to combine the teachings of Ichiriu and Nataraj because the combination of Ichiriu's priority encoder [114] & FLAG circuit [112] and Nataraj's priority encoder/flag logic circuit (215) comprises inputs of match signals and outputs of match addresses and match flag (Ichiriu: figure 1, column 3 line 64-column 4 line 13; Nataraj: figure 4, column 7 line 1-15).

Claim 10:

Ichiriu teaches that a comparand value is stored to the CAM array (101). Ichiriu also teaches that the results of any match between the comparand and a valid CAM word are outputted to a priority encoder (114) for outputting a CAM index (174), such a CAM index is an address of the CAM word corresponding to the selected match signal (column 4 lines 1-8).

Claims 12 and 14:

Ichiriu discloses the invention substantially as claimed. Ichiriu shows (figures 5 and 1) that a CAM array (101) includes a plurality of CAM cells (201) arranged in rows and columns, with each row of CAM cells (201) being coupled to a respective word line and to a respective match line (182). Each CAM (201) includes a compare circuit to compare the content of the memory cell with a comparand signal (figure 5, column 6 line 33-column 7 line 13).

Ichiriu does not explicitly teach a circuit for determining a status of the match line under tested based on result of a search operation and a signal on the match line after confirming proper operation of a control line used to generate a signal.

Nataraj, however, teaches (figure 4) that a match latch circuit (203) comprises latching circuits (221), each latching circuit comprises an AND gate (223) and a latch element (225). Each AND gate (223) has a first input coupled to receive a detect signal (224) and a second input coupled to a corresponding one of the match lines (241). The AND logic gate (223) outputs a logic-level match signal in high or low state according to whether the signal level of the corresponding match line is above or below a threshold. The latch signal (226) is supplied to a latch enable input of the latch element (225) enables the logic-level match signal to pass through the latch element (225) (figure 4, column 48-62).

It would have been obvious to one skilled in the art at the time the invention was made to use Nataraj's match latch circuit (203) having the plurality of AND gates (223) and plurality of latch elements (225) such that each latch element (225) is enabled to

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pass/output a match signal [generated by a AND gate (223)] to Ichiriu's priority encoder (114) and Ichiriu's FLAG circuit (112). One having ordinary skill in the art would be motivated to combine the teachings of Ichiriu and Nataraj because the combination of Ichiriu's priority encoder [114] & FLAG circuit [112] and Nataraj's priority encoder/flag logic circuit (215) comprises inputs of match signals and outputs of match addresses and match flag (Ichiriu: figure 1, column 3 line 64-column 4 line 13; Nataraj: figure 4, column 7 line 1-15).

Claim 13:

Ichiriu also teaches a priority encoder (114), in responsive to the results of any match between the comparand and a valid CAM word, for outputting a CAM index (174), such a CAM index is an address of the CAM word corresponding to the selected match signal (column 4 lines 1-8).

Claim 17:

Ichiriu discloses the invention substantially as claimed. Ichiriu shows (figure 5) that a CAM array (101) includes a plurality of CAM cells (201) arranged in rows and columns, with each row of CAM cells (201) being coupled to a respective word line. Each CAM (201) includes a compare circuit to compare the content of the memory cell with a comparand signal. Each of the match lines (182) is precharged to a high logical level at the start of a comparison operation, but pulled down to a low logical level by the compare circuit within any attached CAM cell (201) that receives comparand signal

which do not match the stored data value. In this configuration, any match line (182) not pulled low constitutes a match signal (figure 5, column 6 lines 33-column 7 line 13).

Ichiriu does not explicitly teach the enabling circuitry for enabling a match line after confirming proper operation of a control line.

Nataraj, however teaches (figure 4) that a match latch circuit (203) comprises latching circuits (221); each latching circuit (221) comprises an AND gate (223) and a latch element (225). Each AND gate (223) has a first input coupled to receive a detect signal (224) and a second input coupled to a corresponding one of the match lines (241). Each of latch elements (225), in response to a latch signal, enables the corresponding logic-level match signal to pass through the latch element (225) after a corresponding AND gate (223), in responsive to a detect signal, outputs a logic-level match signal in high or low state according to whether the signal level of the corresponding match line is above or below a threshold (figure 4, column 48-62).

It would have been obvious to one skilled in the art at the time the invention was made to use Nataraj's match latch circuit (203) having a plurality of AND gates (223) and a plurality of latch elements (225) such that each latch element (225) is enabled to pass/output a match signal [generated by a AND gate (223)] to Ichiriu's priority encoder (114) and Ichiriu's FLAG circuit (112). One having ordinary skill in the art would be motivated to combine the teachings of Ichiriu and Nataraj because the combination of Ichiriu's priority encoder (114) & FLAG circuit (112), and Nataraj's priority encoder/flag logic circuit (215) comprises inputs of match signals and outputs of match addresses

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and match flag (Ichiriu: figure 1, column 3 lines 64-column 4 line 13; Nataraj: figure 4, column 7 line 1-15).

Claim 19:

Ichiriu teaches that each row of CAM cells (201) is coupled to a respective word line (181) and to a respective match line (182) (figure 5, column 6 lines 35-37).

Claims 20-21:

These claims are similar to claims 17 and 19 with additional recited control circuitry for resetting the enabling circuitry. Ichiriu teaches that the CAM array (101) includes circuitry to force validity value with each validity storage cell (202) to a reset state to prevent assertion of a match signal by pulling the match line low for the corresponding row of CAM cells (column 7 lines 49-65, column 8 lines 20-24).

Claim 22:

Ichiriu teaches that a comparand value is stored to the CAM array (101). Ichiriu also teaches that the results of any match between the comparand and a valid CAM word are outputted to a priority encoder (114) for outputting a CAM index (174), such a CAM index is an address of the CAM word corresponding to the selected match signal (column 4 lines 1-8).

Claim 23:

This claim is similar to claim 12 except the additional recited processor and the additional recited circuit for determining a status of a write enable signal to generate a signal on the word lines.

Ichiriu teaches that a system (960) includes a CAM device (961), a CPU (962) and a network processing unit (NPU) (963) (figure 38, column 37 lines 26-45). Ichiriu also teaches an address circuit (103) includes an address selector (125) that responds to the select signal (118) by selecting an address (178) so that an address decoder (127) decodes the selected address (178) to activate one of a plurality of word lines (181) (figures 1 and 2, column 5 lines 17-27).

Claim 30:

Ichiriu discloses the invention substantially as claimed. Ichiriu shows (figure 5) that a CAM array (101) includes a plurality of CAM cells (201) arranged in rows and columns, with each row of CAM cells (201) being coupled to a respective word line. Each CAM (201) includes a compare circuit to compare the content of the memory cell with a comparand signal. Each of the match lines (182) is precharged to a high logical level at the start of a comparison operation, but pulled down to a low logical level by the compare circuit within any attached CAM cell (201) that receives comparand signal which do not match the stored data value. In this configuration, any match line (182) not pulled low constitutes a match signal (figure 5, column 6 lines 33-column 7 line 13).

Ichiriu does not explicitly teach the enabling circuitry for enabling a match line.

Nataraj, however teaches (figure 4) that a match latch circuit (203) comprises latching circuits (221); each latching circuit (221), in responsive to a latch signal and a detect signal, outputs a logic-level match signal in high or low state according to whether the signal level of the corresponding match line is above or below a threshold (figure 4, column 48-62).

It would have been obvious to one skilled in the art at the time the invention was made to use Nataraj's match latch circuit (203) having a plurality of latching circuits (221), each of which to pass/output a match signal to Ichiriu's priority encoder (114) and Ichiriu's FLAG circuit (112). One having ordinary skill in the art would be motivated to combine the teachings of Ichiriu and Nataraj because the combination of Ichiriu's priority encoder (114) & FLAG circuit (112), and Nataraj's priority encoder/flag logic circuit (215) comprises inputs of match signals and outputs of match addresses and match flag (Ichiriu: figure 1, column 3 lines 64-column 4 line 13; Nataraj: figure 4, column 7 line 1-15).

Claims 32:

Claim 32 is rejected for reasons similar to those set forth against claim 19.

Claims 33-34:

These claims are similar to claim 30 with additional recited control circuitry for resetting the enabling circuitry. Ichiriu teaches that the CAM array (101) includes circuitry to force validity value with each validity storage cell (202) to a reset state to

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prevent assertion of a match signal by pulling the match line low for the corresponding row of CAM cells (column 7 lines 49-65, column 8 lines 20-24).

Claim 35:

Ichiriu teaches that a comparand value is stored to the CAM array (101). Ichiriu also teaches that the results of any match between the comparand and a valid CAM word are outputted to a priority encoder (114) for outputting a CAM index (174), such a CAM index is an address of the CAM word corresponding to the selected match signal (column 4 lines 1-8).

Claims 36 and 38:

These claims are similar to claims 30 and 32 except that a router is being recited. Ichiriu teaches a routing device including a CAM device (961) (figure 38, column 37 lines 26-45).

Claim 39:

This claim is similar to claim 33 except that a router is being recited. Ichiriu teaches a routing device including a CAM device (961) (figure 38, column 37 lines 26-45).

Claims 40-41:

Claims 40-41 are rejected for reasons similar to those set forth against claims 34-35.

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Claim 42:

This claim is similar to claims 1 and 12 with additional processing system having a processor. Ichiriu shows a system (960) including a CAM device (961), a CPU (962) and a network processing unit (NPU) (963) (figure 38, column 37 lines 26-45).

Response to Arguments

7. Applicant's arguments filed April 19, 2007 have been fully considered but they are not persuasive.

For claims 12 and 20, the rejection of 112/1st has been withdrawn due to applicant's remark.

For claims 1, 8, 12 and 42, applicant states that the cited references (Ichiriu and Nataraj) fails to teach the feature of confirming proper operation of a control line used to enable output from a match line under test.

Examiner, however, disagrees applicant's remarks. Ichiriu states that at the start of a comparison operation, each of the match lines (182) is precharged to a high logical level (column 6 lines 61-column 7 lines 4). In addition, Nataraj also teaches (figure 4) that within the match latch circuit (203) having a plurality of latching circuits, each latching circuit has a AND gate (223) and a latch element (225). Such a plurality of AND gates (223) and such a plurality of latch elements (225) are connected match lines (241) to used for passing/confirming the match lines (241).

Therefore, it would have been obvious to one skilled in the art at the time the invention was made to use Nataraj's match latch circuit (203) having the plurality of Nataraj's AND gates (223) and plurality of Nataraj's latch elements (225) such that the combination of each AND gate (223) and each latch element (225) is enabled to pass/confirm/output a match signal (as suggested by Nataraj) to Ichiriu's priority encoder (114) and Ichiriu's FLAG circuit (112). One having ordinary skill in the art would be motivated to combine the teachings of Ichiriu and Nataraj because the combination of Ichiriu's priority encoder [114] & FLAG circuit [112] and Nataraj's priority encoder/flag logic circuit (215) comprises inputs of match signals and outputs of match addresses and match flag (Ichiriu: figure 1, column 3 line 64-column 4 line 13; Nataraj: figure 4, column 7 line 1-15).

For claims 17, 20, 23, 30, 33, 36 and 39, applicant also argues that none of the cited reference teach the feature of enabling a match line.

Examiner, however, respectfully traverse applicant's position. Nataraj teaches (figure 4) that within the match latch circuit (203) having a plurality of latching circuits, each latching circuit has a AND gate (223) and a latch element (225). Such a plurality of AND gates (223) and such a plurality of latch elements (225) are connected match lines (241) to used for passing/confirming the match lines (241).

Therefore, it would have been obvious to one skilled in the art at the time the invention was made to use Nataraj's match latch circuit (203) having the plurality of

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Nataraj's AND gates (223) and plurality of Nataraj's latch elements (225) such that the combination of each AND gate (223) and each latch element (225) is enabled to pass/confirm/output a match signal (as suggested by Nataraj) to Ichiriu's priority encoder (114) and Ichiriu's FLAG circuit (112). One having ordinary skill in the art would be motivated to combine the teachings of Ichiriu and Nataraj because the combination of Ichiriu's priority encoder [114] & FLAG circuit [112] and Nataraj's priority encoder/flag logic circuit (215) comprises inputs of match signals and outputs of match addresses and match flag (Ichiriu: figure 1, column 3 line 64-column 4 line 13; Nataraj: figure 4, column 7 line 1-15).

8. **THIS ACTION IS MADE FINAL.** Applicant is reminded of the extension of time policy as set forth in 37 CFR 1.136(a).

A shortened statutory period for reply to this final action is set to expire **THREE MONTHS** from the mailing date of this action. In the event a first reply is filed within **TWO MONTHS** of the mailing date of this final action and the advisory action is not mailed until after the end of the **THREE-MONTH** shortened statutory period, then the shortened statutory period will expire on the date the advisory action is mailed, and any extension fee pursuant to 37 CFR 1.136(a) will be calculated from the mailing date of the advisory action. In no event, however, will the statutory period for reply expire later than **SIX MONTHS** from the mailing date of this final action.

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Any inquiry concerning this communication or earlier communications from the examiner should be directed to Christine T. Tu whose telephone number is (571) 272-3831. The examiner can normally be reached on Mon-Thur. 8:30am-6:00pm.

If attempts to reach the examiner by telephone are unsuccessful, the examiner's supervisor, Jacques Louis-Jacques can be reached on (571) 272-6962. The fax phone number for the organization where this application or proceeding is assigned is 571-273-8300.

Information regarding the status of an application may be obtained from the Patent Application Information Retrieval (PAIR) system. Status information for published applications may be obtained from either Private PAIR or Public PAIR. Status information for unpublished applications is available through Private PAIR only. For more information about the PAIR system, see <http://pair-direct.uspto.gov>. Should you have questions on access to the Private PAIR system, contact the Electronic Business Center (EBC) at 866-217-9197 (toll-free). If you would like assistance from a USPTO Customer Service Representative or access to the automated information system, call 800-786-9199 (IN USA OR CANADA) or 571-272-1000.



Christine T. Tu
Primary Examiner
Art Unit 2117

July 9, 2007